

Biometric Passports

Security and Privacy

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What is a biometric passport?

- ▶ A passport with an embedded chip containing biometric data
- ▶ Also known as an **e-passport** or **digital passport**
- ▶ Standardised by the International Civil Aviation Organisation (ICAO) as Document 9303 [1]
- ▶ Identified by this symbol: 

What data does it contain?

All biometric passports contain a chip with the following data:

- ▶ A copy of the information printed on the passport
- ▶ A digital photo of the passport holder
- ▶ Digital signatures from the issuing country

Issuing countries may optionally include additional data:

- ▶ Fingerprints
- ▶ Iris scans
- ▶ Additional personal information (e.g. address)
- ▶ Additional document information (e.g. observations)

A Brief History

- ▶ 1968: Development begins
- ▶ 1980: Machine-readable passport first standardised
- ▶ 1988: Work begins on biometric systems
- ▶ 1998: First biometric passport issued by Malaysia [2]
- ▶ 2006: Biometric passport standardised



Live Demo

- ▶ Reading the passport chip with an NFC-enabled phone
- ▶ Written in TypeScript (big mistake)



Basic Access Control (BAC)

- ▶ The data in the passport chip cannot be accessed without first establishing a secure channel
 - ▶ One way of doing this is with **Basic Access Control** (BAC)
 - ▶ We will call the scanner the **interface device** (IFD) and the chip the **integrated circuit** (IC)
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- ▶ The hash function **H** is SHA-1
 - ▶ The encryption function **E** is Triple-DES
 - ▶ The MAC function **M** is ISO 9797-1 MAC algorithm 3

Step 0: Document Access Key Calculation

P = passport number

D = date of birth

E = date of expiry

$MRZ = P \parallel D \parallel E$

$K_{seed} = H(MRZ)_{0,\dots,15}$

$KDF(K, c) = H(K \parallel c)_{0,\dots,15}$

$K_{Enc} = KDF(K_{seed}, 1)$

$K_{Mac} = KDF(K_{seed}, 2)$

\parallel denotes concatenation

$X_{a,\dots,b}$ denotes the bytes a to b of X

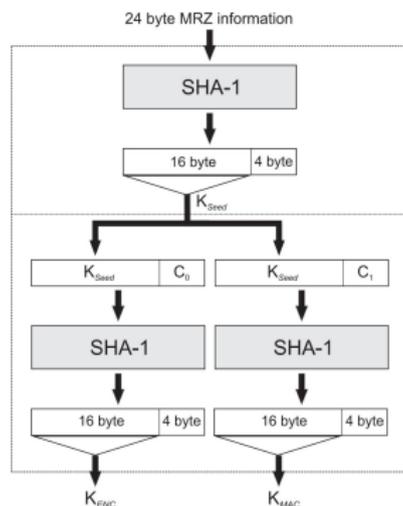


Figure 1: BAC key derivation [4]

Step 1: Basic Access Control

Inspection System (IFD)

$$\text{RND}_{\text{IFD}}, K_{\text{IFD}} \in \{0, 1\}^{64}$$

$$X \leftarrow \text{RND}_{\text{IFD}} \parallel \text{RND}_{\text{IC}} \parallel K_{\text{IFD}}$$

$$E_{\text{IFD}} \leftarrow \mathbf{E}_{K_{\text{Enc}}}(X)$$

$$M_{\text{IFD}} \leftarrow \mathbf{M}_{K_{\text{Mac}}}(E_{\text{IFD}})$$

$$\xleftarrow{\text{RND}_{\text{IC}}}$$

$$\xrightarrow{E_{\text{IFD}} \parallel M_{\text{IFD}}}$$

Decrypt and verify $E_{\text{IC}} \parallel M_{\text{IC}}$

$$\text{KS}_{\text{Seed}} = K_{\text{IFD}} \oplus K_{\text{IC}}$$

$$\xleftarrow{E_{\text{IC}} \parallel M_{\text{IC}}}$$

Passport (IC)

$$\text{RND}_{\text{IC}} \in \{0, 1\}^{64}$$

Decrypt and verify $E_{\text{IFD}} \parallel M_{\text{IFD}}$

$$K_{\text{IC}} \in \{0, 1\}^{128}$$

$$Y \leftarrow \text{RND}_{\text{IC}} \parallel \text{RND}_{\text{IFD}} \parallel K_{\text{IC}}$$

$$E_{\text{IC}} \leftarrow \mathbf{E}_{K_{\text{Enc}}}(Y)$$

$$M_{\text{IC}} \leftarrow \mathbf{M}_{K_{\text{Mac}}}(E_{\text{IC}})$$

$$\text{KS}_{\text{Seed}} = K_{\text{IFD}} \oplus K_{\text{IC}}$$

Step 2: Session Key Derivation

Once the shared secret $KS_{Seed} = K_{IFD} \oplus K_{IC}$ is established:

$$KS_{Enc} = \mathbf{KDF}(KS_{Seed}, 1)$$

$$KS_{Mac} = \mathbf{KDF}(KS_{Seed}, 2)$$

The **send sequence counter** (SSC) is further initialised:

$$SSC = (RND_{IC})_{4,\dots,7} \parallel (RND_{IFD})_{4,\dots,7}$$

Step 3: Secure Messaging

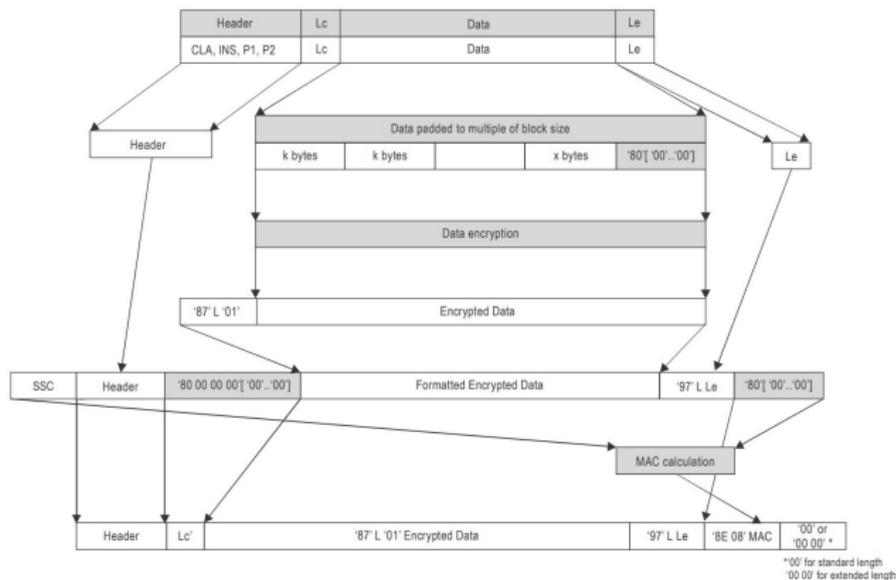


Figure 2: SM command APDU structure [1]

Passive Authentication

With Secure Messaging enabled, the IFD can read the holder's data from the IC and verify its authenticity and integrity using **Passive Authentication**.

- ▶ Each issuing country has a **Country Signing Certificate Authority** (CSCA) which issues certificates for national **Document Signers** (DS)
- ▶ CSCA certificates must be acquired from a trustworthy source (e.g. ICAO PKD)
- ▶ The **Document Security Object** (SO_D) contains digital signatures over hashes of the data in the passport, as well as the DS certificate
- ▶ The inspection system builds and validates a certificate chain from a **Trust Anchor** to the DS certificate
- ▶ The inspection system finally verifies the digital signatures in SO_D

Attacks

- ▶ **E-Passport: Cracking Basic Access Control Keys**, Liu et al. (2007)
Due to the low entropy of the MRZ, the BAC keys can be cracked in a matter of seconds using specialised hardware.
- ▶ **A Traceability Attack Against e-Passports**, Chothia et al. (2010)
By measuring the time taken for the IC to respond, the movements of an individual passport can be traced.
- ▶ **ePassport: Side Channel in the Basic Access Control**, Sportiello et al. (2014)
Further timing analysis can be used to recover the MRZ without eavesdropping.

E-Passport: Cracking Basic Access Control Keys [4]

- ▶ Eavesdrop RND_{IC} , $E_{IFD} \parallel M_{IFD}$, $E_{IC} \parallel M_{IC}$ and the entire subsequent communication C
- ▶ Run a key search on the MRZ information to match the most significant eight bytes of E_{IC} .
- ▶ C can then be decrypted.

This becomes feasible if the MRZ has low entropy, due to:

- ▶ Passport numbers that are sequential, structured, include a checksum, or are otherwise predictable
- ▶ Passport expiry dates having a small range
- ▶ An attacker being able to narrow down the date of birth

E-Passport: Cracking Basic Access Control Keys [4]

- ▶ At the time, a lot of passports had these issues!
- ▶ **Germany**: 4 digits for local authority (of which there are 295), remaining 5 digits sequential
- ▶ **Netherlands**: Begins with fixed character “N”, ends with a check digit, remaining 7 digits sequential

The authors demonstrated that, with just a photo of the passport holder and some knowledge about the dependency between passport numbers and expiry dates, the keys for a German passport could be cracked in ≈ 22 seconds and a Dutch passport in ≈ 10.3 seconds.

A Traceability Attack Against e-Passports [5]

- ▶ Eavesdrop $E_{IFD} || M_{IFD}$ during a legitimate BAC exchange
- ▶ Initiate a BAC exchange with the IC and send it random data
- ▶ Initiate a BAC exchange with the IC and send it the recorded message
- ▶ Compare the time taken for the IC to respond

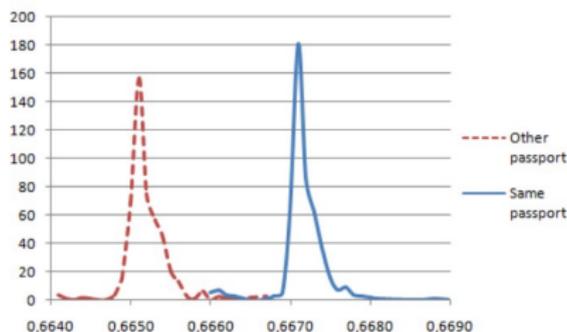


Figure 3: Sampled response times (British passport) [5]

A Traceability Attack Against e-Passports [5]

- ▶ IC checks the MAC before comparing nonces
- ▶ If the MAC is invalid (i.e. K_{Mac} is wrong), the IC will not compare nonces, and the response time will be shorter
- ▶ Could be used to, for example, build a bomb that detonates in the presence of a specific person [6]
- ▶ All 10 passports from 6 countries that the authors tested were vulnerable
- ▶ The French passport was even worse: it went against the spec and explicitly gave a different error code!

ePassport: Side Channel in the Basic Access Control [7]

- ▶ Set a generic E_{IFD} and a random M_{IFD} .
- ▶ Vary the i th byte of M_{IFD} and measure the time taken for the IC to respond.
- ▶ The value that causes the longest mean response time is the correct value of the i th byte.
- ▶ Work through the bytes to find the valid MAC for E_{IFD} .
- ▶ Use specialised hardware to crack the MRZ.

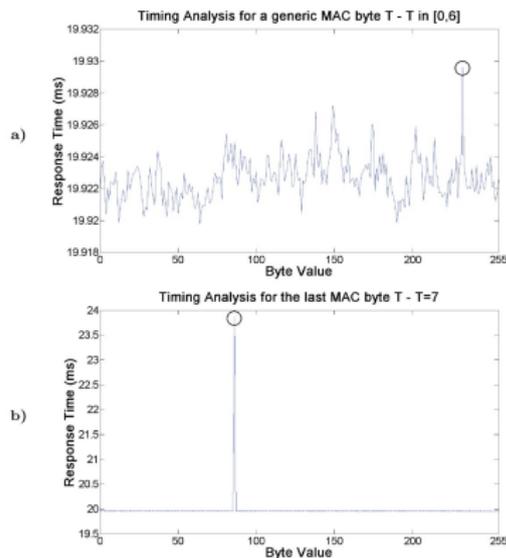


Figure 4: Timing analysis [7]

ePassport: Side Channel in the Basic Access Control [7]

- ▶ Requires an interaction of ≈ 85 hours with the IC
- ▶ The authors suggest that this could be achieved by performing the attack from malware on the target's phone:
 - ▶ Phone and passport are kept in the same pocket or bag
 - ▶ Malware uses the phone's NFC reader to communicate with the passport
 - ▶ Once the keys are cracked, the malware can extract further data from the passport and send it to the attacker
 - ▶ Alternatively the passport could be used remotely through a relay attack

Mitigations

- ▶ Ensure that the IC does not leak information through the timing of its responses
- ▶ Improve the entropy of the MRZ (e.g. by using a random passport number)
- ▶ Shield the chip so it can only be read if the passport is open
 - ▶ The US passport has a thin metal mesh in the cover [8]
- ▶ What to do about the fundamentally insecure DES encryption?

Password Authenticated Connection Establishment (PACE)

- ▶ PACE is an alternative to BAC that uses asymmetric cryptography
- ▶ Generates strong session keys independent of the strength of the password (in this case, the MRZ)
- ▶ Based on Diffie-Hellman key exchange
- ▶ Further communication is encrypted using AES
- ▶ See ICAO Doc 9303 Part 11, Section 4.4 [1]

Cryptographic Weaknesses

- ▶ EU passports since 2014 have been required to use PACE [9]
- ▶ Until 2018 it was still required to support BAC [1]
- ▶ PACE prevents eavesdropping but if BAC is still supported, the MRZ can still be cracked

LDS2 Applications

- ▶ In 2021, the ICAO introduced **Logical Data Structure 2** (LDS2), which allows for additional data to be stored
- ▶ This includes travel history, visa records, and additional biometric data

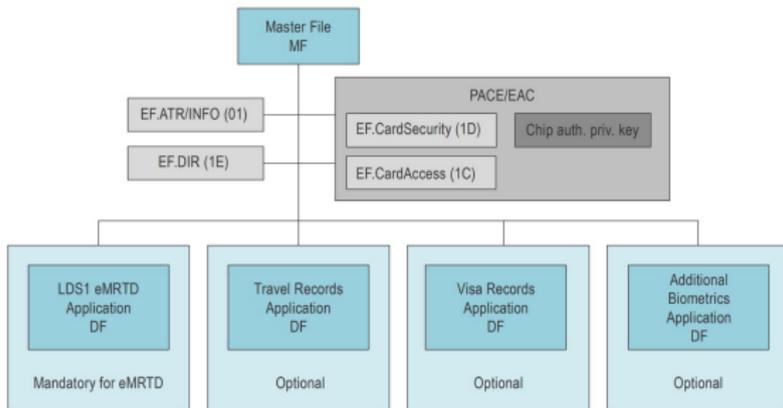


Figure 5: Applications for LDS1 and LDS2 [1]

Mandatory recording of travel history

- ▶ Up to the issuing country whether to enable this feature
- ▶ If enabled, it is mandatory for all countries to record the holder's travel history
- ▶ Currently, countries that don't want visitors stigmatised can issue visas and record entries and exits on separate documents

Tag	Tag	Content	Mandatory /OPTIONAL	Format	Example
'SF44'		Embarkation/Debarcation State (copy for SEARCH RECORD)	M	F (3) A	USA
'73'	Entry / Exit Travel Record (signed info)				
	'SF44'	Embarkation/Debarcation State	M	F (3) A	USA
	'SF4C'	Visa approvals, refusals, and revocations	O	V (50) A,N,S,U	Free-form text
	'SF45'	Travel date (Date of entry/exit)	M	F (8) N	20120814 (yyyymmdd)
	'SF4B'	Inspection authority	M	V (10) A,N,S	CBP
	'SF46'	Inspection location (Port of Entry/Exit)	M	V (10) A,N,S	SFO
	'SF4A'	Inspector reference	M	V (20) A,N,S	SFO00001234
	'SF4D'	Result of inspection	O	V (50) A,N,S,U	Free-form text
	'SF49'	Mode of travel	O	F (1) A	A (Air), S (Sea), L (Land)
	'SF48'	Duration of stay (days)	O	V (2) B	'00FF' (255 days)
	'SF4E'	Conditions holder is required to observe while in the issuing State	O	V(50) A,N,S,U	Free-form text
'SF37'		Authenticity token (Signature)	M	V (140) B	'SF' '37' Len (Signature)
'SF38'		Reference (record number) to LDS2-TS Signer certificate in Certificates Store	M	F (1) B	'01'...'FE'

Figure 6: Entry/Exit Record [1]

Irremovable messages

- ▶ Countries can add irremovable messages to the passport
- ▶ For example, a country could mark a person as “suspicious” with no reason given and no way to remove it
- ▶ This would make it difficult for the holder to travel [10]

Conclusion

- ▶ Biometric passports allow for much faster border control without compromising security
- ▶ Most of the security issues are fixed and just need time to be rolled out
- ▶ The privacy issues are just getting started...

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